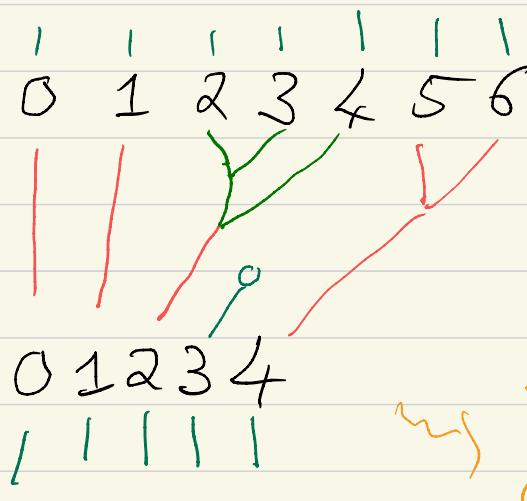
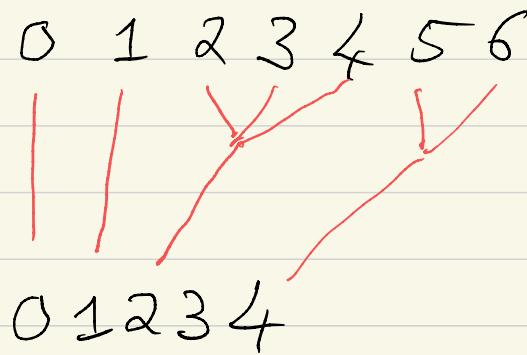


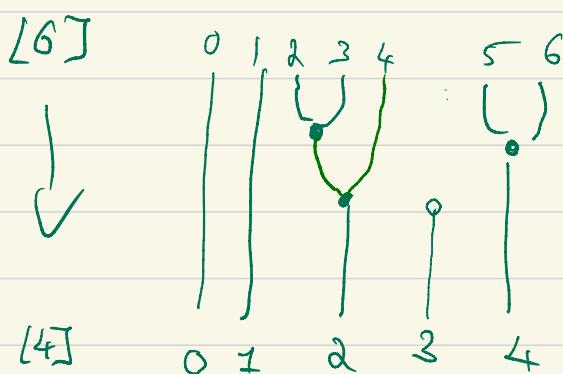
lecture 2 Simplicial sets

The category Δ has as objects the finite ordinals $[n] = \{0 < 1 \dots < n\}$
 (in part. $[0] = \{0\}$)

Morphisms are (weakly) monotone functions:



two "graphical" generators



(cf. μ or η of a monad)

Presenting Δ by generators and relations

$$d^i : \mathbf{n} - 1 \rightarrow \mathbf{n} \quad 0 \leq i \leq n \quad (\text{cofaces})$$

$$s^j : \mathbf{n} + 1 \rightarrow \mathbf{n} \quad 0 \leq j \leq n \quad (\text{codegeneracies})$$

$$d^i = \left| \begin{array}{c|c|c|c|c} & 0 & i-1 & i & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right| \quad s^j = \left| \begin{array}{c|c|c|c|c} & 0 & j-2 & j & j+1 & j+2 & n+1 \\ \hline 0 & | & | & | & | & | & | \\ & | & | & | & | & | & | \\ & | & | & | & | & | & | \\ & | & | & | & | & | & | \end{array} \right|$$

$$\left\{ \begin{array}{l} d^i d^j = d^{j+1} d^i \quad (i \leq j) \quad (1) \\ d^i s^j = s^{j+1} d^i \quad (i \leq j) \quad (2) \\ s^i d^j = 1 = s^j d^{j+1} \quad (a), (b) \\ s^i d^j = d^{j+1} s^i \quad (i < j-1) \quad (3) \\ s^i s^j = s^{j+1} s^i \quad (i < j) \quad (f)(i < j-1), (c) \quad (i = j-1) \end{array} \right.$$

(1)-(4) are "naturality conditions" = "pulling strings differently"

(a)-(c) are the monad / monoid laws

$$\left| \begin{array}{c|c|c|c|c} & 0 & i & j & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right| \stackrel{(1)}{=} \left| \begin{array}{c|c|c|c|c} & 0 & i & j+1 & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right|$$

$$\left| \begin{array}{c|c|c|c|c} & 0 & i & j & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right| \stackrel{(2)}{=} \left| \begin{array}{c|c|c|c|c} & 0 & i & j & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right|$$

$$\left| \begin{array}{c|c|c|c|c} & 0 & i & j & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right| \stackrel{(3)}{=} \left| \begin{array}{c|c|c|c|c} & 0 & i & j-1 & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right|$$

$$\left| \begin{array}{c|c|c|c|c} & 0 & i & j & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right| \stackrel{(4)}{=} \left| \begin{array}{c|c|c|c|c} & 0 & i & j-1 & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right|$$

$$\left| \begin{array}{c|c|c|c|c} & 0 & i & j & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right| \stackrel{(a)}{=} \left| \begin{array}{c|c|c|c|c} & 0 & i & j & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right|$$

$$\left| \begin{array}{c|c|c|c|c} & 0 & i & j & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right| \stackrel{(b)}{=} \left| \begin{array}{c|c|c|c|c} & 0 & i & j & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right|$$

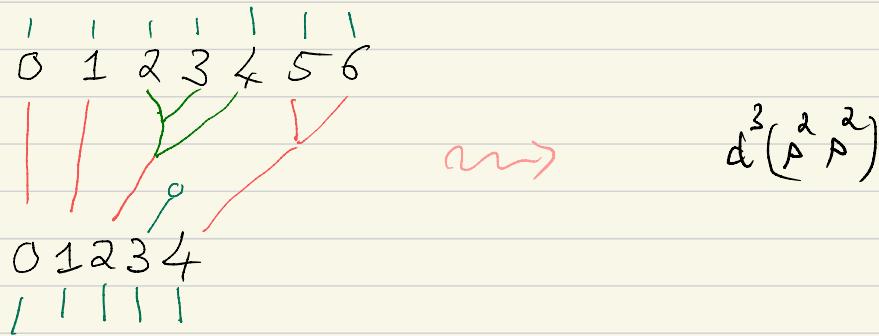
$$\left| \begin{array}{c|c|c|c|c} & 0 & i & j & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right| \stackrel{(c)}{=} \left| \begin{array}{c|c|c|c|c} & 0 & i & j & n-1 \\ \hline 0 & | & | & | & | \\ & | & | & | & | \\ & | & | & | & | \end{array} \right|$$

Proving that this presentation is indeed a presentation of \mathbb{D}

- Orient all equations from left to right and prove that they form a confluent and terminating rewriting system \rightarrow uniqueness of normal forms
(See next page)

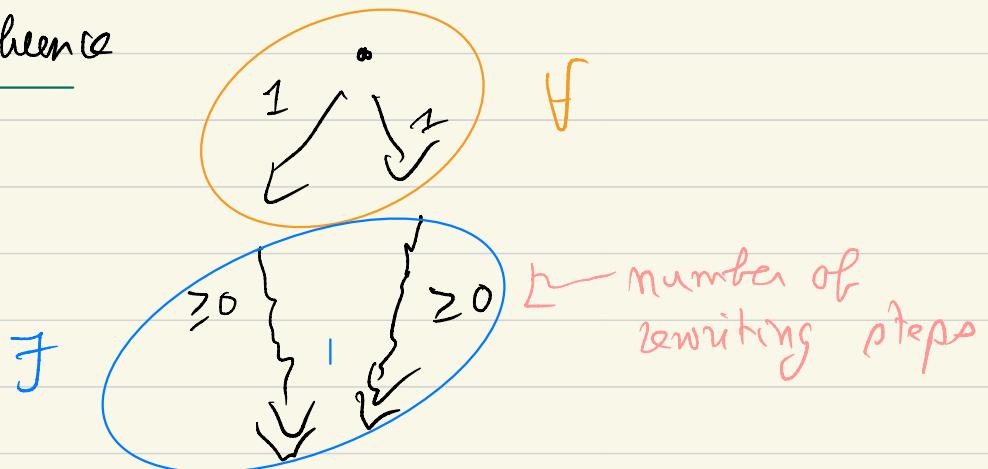
- Observe that the normal forms are of the form
$$u_2 \dots u_p, \text{ together with a strictly decreasing map } [1, p] \xrightarrow{i} [n]$$
for some n
s.t. for all j $u_j = d^i \circ \dots \circ d^i$
$$\text{or } u_j = \underbrace{d^i \circ \dots \circ d^i}_{q \text{ times}}, \text{ for some } q$$

- Observe that these normal forms encode precisely the descriptions p.1:

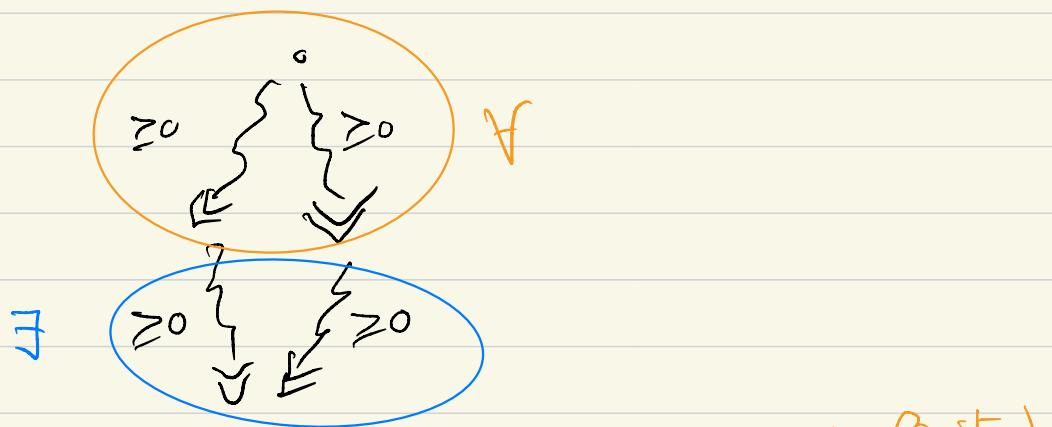


Rewriting theory background

- Local confluence



- Confluence:



- strong normalization $\not\vdash \cdot \rightarrow \cdot \rightarrow \cdots \rightarrow \cdots$

- normal form: $\bullet \not\rightarrow$

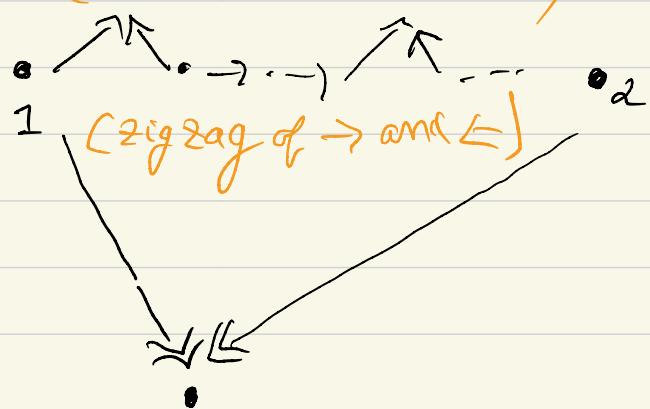
- Local confluence + strong normalization \Rightarrow confluence

(Newman Lemma)

If confluence +
strong norm., then

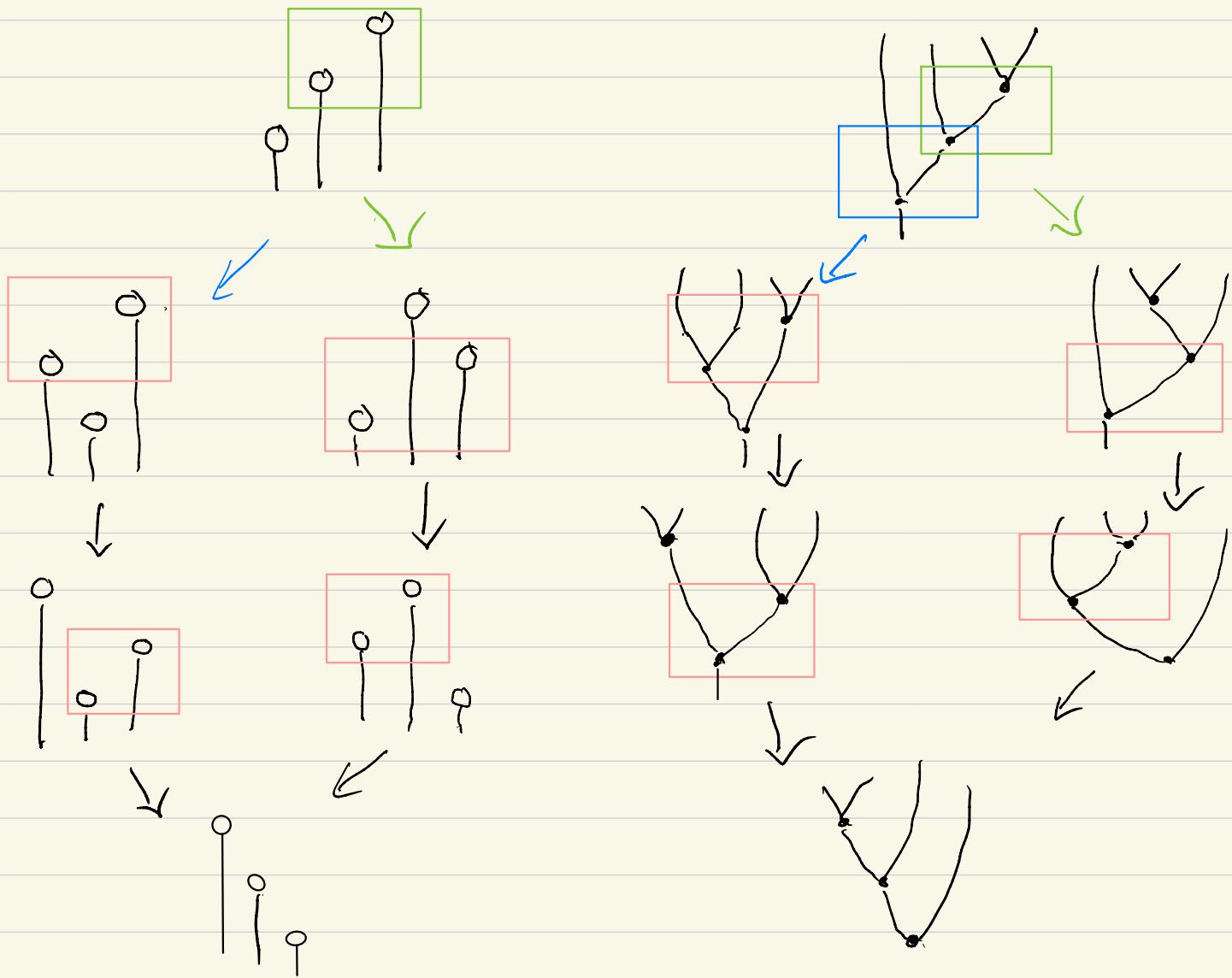
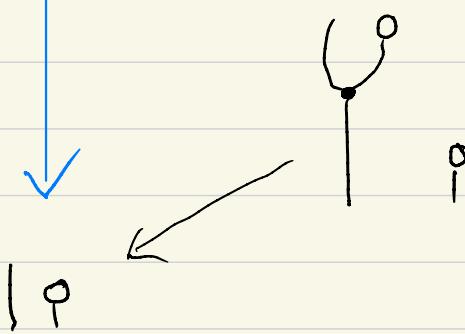
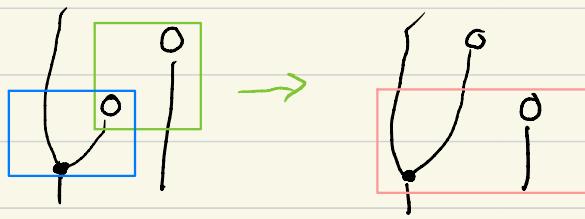
$\Rightarrow \bullet_1$ and \bullet_2 have the same

normal form.



Local confluence checking

For example



Simplicial Sets

More generally, a simplicial object in a category C is a functor $\Delta^{\text{op}} \rightarrow C$

A simplicial set is a presheaf $X: \Delta^{\text{op}} \rightarrow \text{Set}$

The category of simplicial sets is $\text{Set}_{\Delta} = \text{Set}^{\Delta^{\text{op}}}$

↑ also written $\widehat{\Delta}$

Terminology: An element of X_n is called an element of X over n , or an n -simplex

Notation: we write $d_i x$ for $X d_i x$ and $\delta_i x$ for $X \delta_i x$

Hence $d_i: X_n \rightarrow X_{n-1}$ and $\delta_i: X_n \rightarrow X_{n+1}$

Drawing the elements of a simplicial set: *Anticipating topological realisation!*

at X_0

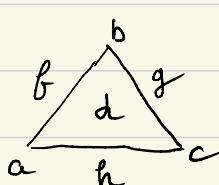
$f \in X_1$

$d \in X_2$

$A \in X_3$ tetrahedron

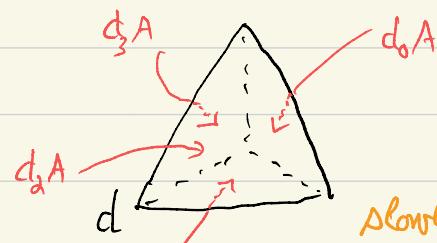
i

$a \xrightarrow{i} b$



$$\begin{aligned} \text{where } a &= d_1 b \\ b &= d_0 f \end{aligned}$$

$$\begin{aligned} \text{where } f &= d_2 d \\ g &= d_0 d \\ h &= d_2 d \\ a &= d_1 f = d_1 h \end{aligned}$$



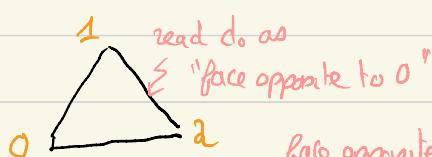
Slowly next page

\Downarrow
a decoration of

\Downarrow
a decoration of

\Downarrow
a decoration of

$0 \longrightarrow 1$



read d_0 as
"face opposite to 0"



face opposite to 3

$b = d_{(0,1,2)} a$

$d = d_{(1,2,3)} a$

Thus (repeated) d_i 's take iterated faces.

⇒ injections

$X_f x \in X_m$

$f: [m] \rightarrow [n]$

$x \in X_n$

result of transporting x along f .

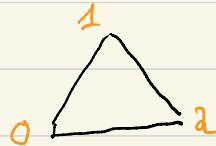
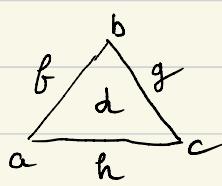
A handy notation for iterated faces: for $f: [m] \rightarrow [n]$ injective we write

$$X_f x = d_{[n]/f([m])} x$$

↑

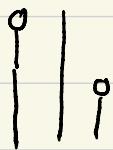
$d_I x = (\text{decoration of})$
the face obtained by intersecting the codimension 1 faces of x opposite to $i \in I$

Combinatorics of factoring faces

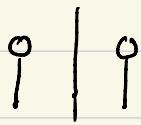


We have

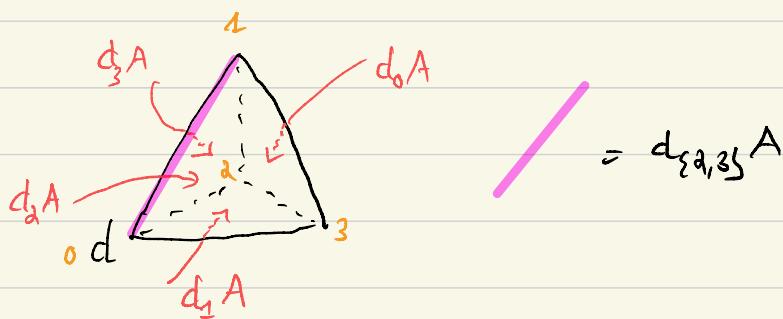
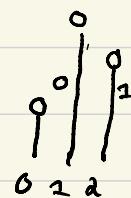
$$b = d_0 f = d_0 d_2 a$$



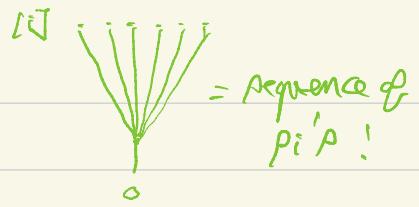
$$b = d_{\{0,2\}} a$$



$$b = d_1 g = d_1 d_0 a$$

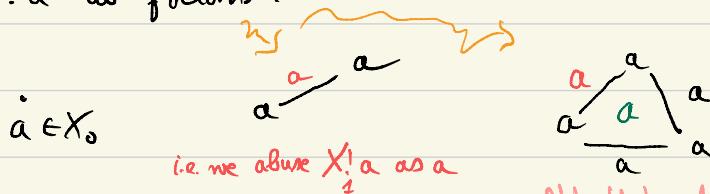


Degeneracies



What about π_i ? Let $a \in X_0$ and let $!_i$ be the unique map $[i] \rightarrow [0]$.

Then we draw $X!_i a$ as follows:

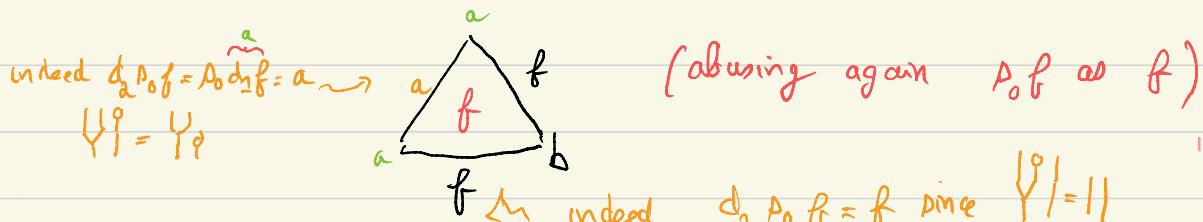


Note that abuses are consistent:

$$d_2^2 a = d_2 X!_2 a = X(d_2^2) a = X!_2 a = a$$

Other example: let $f \in X_1$ then we draw $\rho_0 f \in X_2$ as follows:

(character for ρ_i : glueing copies of f along their i -th face)

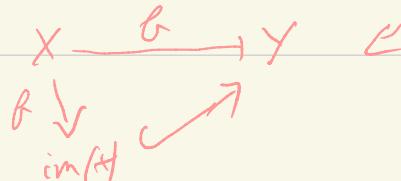


Definition An n -simplex x is called degenerate if there exists a surjection ($\neq \text{id}$) $f: [n] \rightarrow [m]$ (hence $m \leq n$) and an m -simplex y p.t. $x = Xfy$.

(The above examples are degenerate simplices)

Exercise Prove that one can replace "surjection $f \neq \text{id}$ " by "any $f: [n] \rightarrow [m]$ p.t. $m < n$ " (hint: use the epi-mono factorisation in Δ , showing on the way that it exists and is inherited from the one in Set).

Exercise Show that if x is non-degenerate, and f, y are such that $x = Xfy$, then f is injective.
(hint: use again the epi-mono factorisation)



The standard simplex Δ^n

Δ^n is the simplicial set $\mathcal{Y}_n \quad (\mathcal{Y}: \Delta \rightarrow \widehat{\Delta})$

Explicitly: $\binom{\Delta^n}{m} = \Delta(m, n)$

Exercise Show that the non-degenerate simplices of Δ^n are exactly the injective maps.

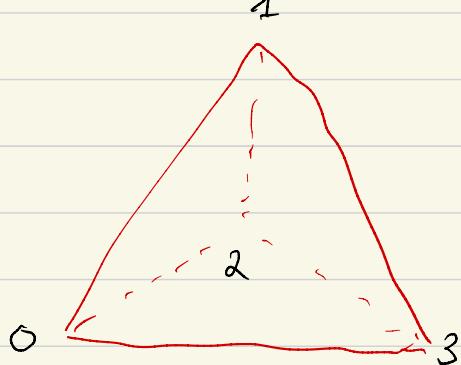
Hint: note that $f: [k] \rightarrow [n]$ is non-injective iff

$$[k] \xrightarrow{f} [n]$$

epi + $k' \subset k$ ↗
 [k'] ↗
 mono

Therefore, say, Δ^3 has a unique non-degenerate 3-simplex (increasing injections = identities), justifying to draw it

as



(iterated)

Its faces are given by the subsets of $\{0, 1, 2, 3\}$

(see also simplicial complex below)

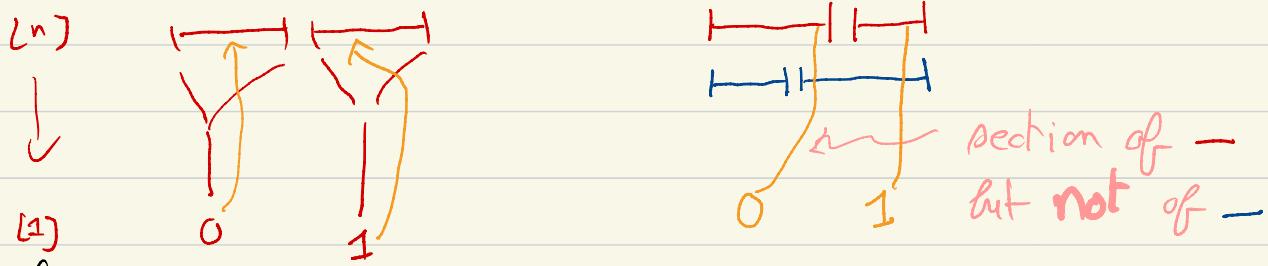
Eilenberg - Zilber's lemma

Proposition Let $X \in \Delta$ and $x \in X_n$. Then there exists a unique pair (f, g) s.t. • $f: [n] \rightarrow [m]$ is surjective
 • $y \in X_m$ is non-degenerate
 • $X f y = x$

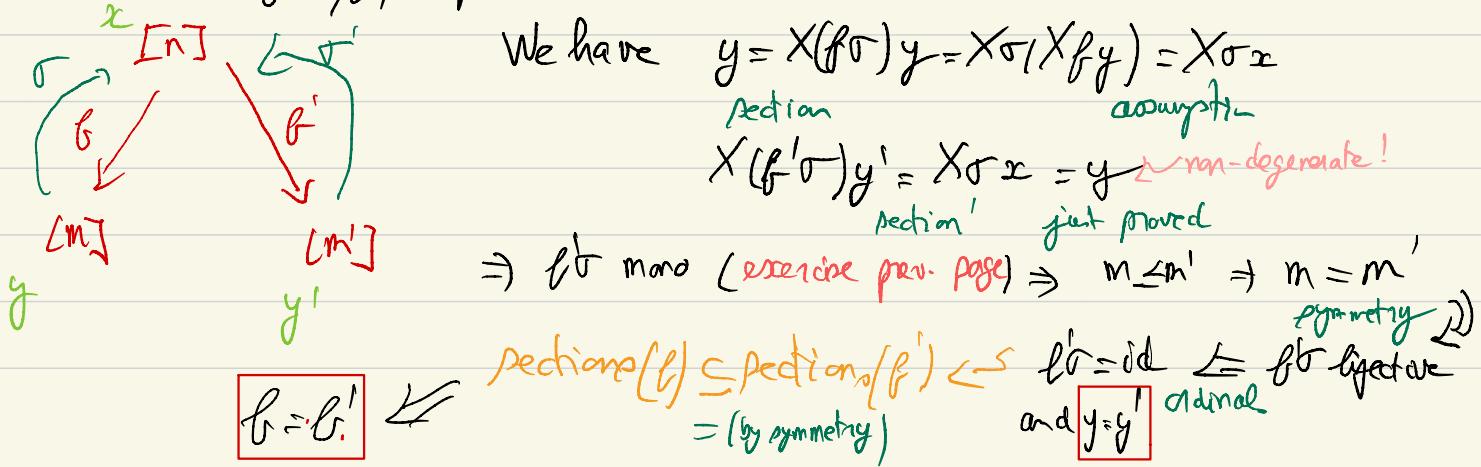
Proof Existence: If x is non-degenerate, take $f = \text{id}$, $y = x$.
 If x is degenerate, then we apply iteratively the definition of being degenerate (induction on n).

Uniqueness: We use easy properties of Δ : • if $f: [m] \rightarrow [n]$ is surjective, then it admits a section (any set-theoretic section is automatically monotone):

• Moreover, f is characterized by its set of sections



Suppose (f, g) and (f', g') satisfy the assumptions. Take σ, σ' sections of f, f' , resp.:



Simplicial complexes

Simplicial complexes are simpler than simplicial sets, but less flexible.

A simplicial complex is a pair (V, C) , where

V is a set of vertices (the 0-simplices), and where

$$C \subseteq \binom{V}{\leq n}$$

, and

- $x \in C \Rightarrow x \neq \emptyset$
- $\{v_1, v_2\} \subseteq C$
- $\emptyset \neq y \subseteq x \in C \Rightarrow y \in C$

The elements of C of cardinal n are the n -simplices.

is $\left(\begin{array}{c} \text{circle} \\ \text{as} \\ \vdots \end{array}\right) g$

\vdots

__ simplicial complex

simplicial set

$$V = \{v_0, v_1, v_2\}$$

$$C = \{\{v_0, v_1, v_2\}, \{v_0, v_1\}, \{v_1, v_2\}, \{v_0, v_2\}\}$$

A morphism $\tilde{\epsilon}: (V, C) \rightarrow (V', C')$ of simplicial complexes is a function $\tilde{\epsilon}: V \rightarrow V'$ such that $\forall A \in C \quad \tilde{\epsilon}(A) \in C'$.

Relating simplicial complexes and simplicial sets

We make a little twist on both sides to establish a reflection.

- We consider ordered simplicial complexes, which are (V, C) + a total order on V . This affects the morphisms only! we accept only those $\Xi : (V, C) \rightarrow (V', C')$ such that $\Xi : V \rightarrow V'$ is monotone (= non-decreasing).
- We consider 0-simplices - ordered simplicial sets, which are X + a total order on X_0 . Again, this affects the morphisms only: we accept only those $f : X \rightarrow Y$ such that $f_0 : X_0 \rightarrow Y_0$ is monotone.

We denote these two categories by Spcx and Sset .

We write TO for the category of totally ordered sets and monotone functions.

From Spcx to Sset . We define, for $(V, C) \in \text{Spcx}$ \sim_V if $\{v_1, v_2\} \in C$, then $v_1 \sim v_2 \in (X_C)_2$

$$(X_C)_n = \{\varphi \in \text{TO}[m], V] \mid \varphi([m]) \in C\} \quad (\text{note that } (X_C)_0 \text{ is naturally ordered})$$

For $f : [m] \rightarrow [n]$, we set $(X_C)_f : (X_C)_n \rightarrow (X_C)_m = \varphi \mapsto \varphi \circ f$

(note that $(\varphi \circ f)[m] = \varphi[f[m]] \subseteq \varphi[n] \in C$, hence $(\varphi \circ f)[m] \in C$)

In English, simplices are lists $(v_0 \dots v_n)$ where there may be repetitions = degenerate simplices!

For $d^i = (0 \dots \hat{i} \dots m+1) : (v_0 \dots v_{m+1}) \mapsto (v_0 \dots \hat{v}_i \dots v_{m+1})$

For $s^i = (0, \dots, i, \dots, m) : (v_0 \dots v_m) \mapsto (v_0 \dots v_i, v_i \dots v_m)$

For $\Xi : (V, C) \rightarrow (V', C')$ we set $(\Xi^*)_n \varphi = \Xi \circ \varphi$ (note: $\Xi(\varphi[n]) \in C'$)

Thus $(V, C) \mapsto X_C$, $\Xi \mapsto \Xi^*$ defines a functor

$\text{Spcx-to-Sset} : \text{Spcx} \rightarrow \text{Sset}$

From $\underline{\text{Sset}}$ to $\underline{\text{Spc}_\infty}$. Let X be a simplicial set. We define (V_X, \mathcal{L}_X) as follows:

$$V_X = X_0 \quad (\text{with its total order})$$

$$\mathcal{L}_X = \{ \text{Vert}(n, x) \mid x \in X_n \} = \text{Vert}(\text{elt}(X)), \text{ where}$$

$$\text{Vert}(n, x) = \{ X f x \mid f: [0] \rightarrow [n] \}$$

Claim: \mathcal{L}_X is closed under subsets.

Proof: We introduce notation: $\begin{cases} \text{if } f: [0] \rightarrow [n] \text{ then } \underline{f} = f|0 \\ \text{if } j \in [n], \text{ then } \bar{j}: [0] \rightarrow [n] = 0 \mapsto j \end{cases}$

Let $A \subseteq \text{Vert}(n, X)$. Let $A' = \{ j \mid X \bar{j} x \in A \}$

Let $\iota: [m] \rightarrow [n]$ be s.t. ι is injective and its image is A' .

$$\text{Then } \text{Vert}(m, X \iota x) = \{ X g(X \iota x) \mid g: [0] \rightarrow [m] \}$$

$\llcorner X(\underline{g}) x$

$$= \{ X f x \mid f \text{ factors through } \iota \}$$

$$= \{ X f x \mid \underline{f} \in A' \} = \{ X f x \mid X \bar{f} x \in A \} = A$$

Let $u: X \rightarrow Y$. Then we have $u_0: X_0 \rightarrow Y_0$.

Claim: u_0 is a morphism. Indeed, we have $\underline{u_0}(\text{Vert}(n, x)) \subseteq \text{Vert}(n, u(x))$:

$$\begin{array}{ccc} x \in X_n & \longrightarrow & Y_n \ni u_n(x) \\ \downarrow X_f & & \downarrow Y_g \\ X_0 & \xrightarrow{u_0} & Y_0 \end{array}$$

\nwarrow

$\exists f, x \nearrow$

$\forall f, x \nearrow$

$\exists g, u(x) \nearrow$

Thus we have a functor $\underline{\text{Sset}} \rightarrow \underline{\text{Spc}_\infty}$: $X \mapsto (X_0, \mathcal{L}_X)$, $u \mapsto u_0$.

Theorem The functors SSet -to- Spc and Spc -to- SSet form a reflection, i.e. SSet -to- Spc : $\text{SSet} \xrightarrow{\perp} \text{Spc}$: Spc -to- SSet

and Spc -to- SSet is full and faithful.

Proof Full and faithfulness follows from

$$\underline{\Sigma}(v) = \underline{\Sigma} \circ \bar{v} = \underline{\Sigma}^* \bar{v}$$

i.e. we can retrieve $\underline{\Sigma}$ from Spc -to- $\text{SSet}(\underline{\Sigma}) = \underline{\Sigma}^*$.

For X and (\mathcal{V}, V) , we need to establish

$$\underline{\Sigma}: X_0 \rightarrow \mathcal{V} \quad \text{s.t. } \forall (n, x) \in \text{el}(X) \quad \underline{\Sigma}(\text{Vert}(n, x)) \in \mathcal{C}$$

$$\underline{\Sigma}_n: X_n \rightarrow (\mathcal{C})_n \quad \text{natural}$$

- Given $\underline{\Sigma}$, define $\underline{\Sigma}(x_0) = \underline{\Sigma}_0(x_0)$. We have to check $\underline{\Sigma}(\text{Vert}(n, x)) \in \mathcal{C}$, which follows from the following claim

Claim $\underline{\Sigma}(\text{Vert}(n, x)) \subseteq (\underline{\Sigma}_n x)[n]$:

$\in \mathcal{C}$ by def. of \mathcal{C}

$$\begin{array}{ccc} x & X_n & \xrightarrow{\underline{\Sigma}_n} (\mathcal{C})_n \\ \downarrow f & \downarrow & \downarrow -of \\ x_{f(x)} & X_0 & \xrightarrow{\underline{\Sigma}_0} (\mathcal{C})_0 = \text{To}[\{x\}], \mathcal{V} \\ & \text{or} & \text{more precisely} \\ & \underline{\Sigma}_n(x) \circ f & \underline{\Sigma}_0 = \underline{\Sigma}_n(x)[n] \end{array}$$

- Given $\underline{\Sigma}$, (a) the bijection forces $\underline{\Sigma}_0: \underline{\Sigma}_0(x_0) = \underline{\Sigma}(x_0)$

(b) Claim Any $\underline{\Sigma}$ is determined by its $\underline{\Sigma}_0$ component: Choose n and $i \in [n]$: put $f = \bar{i}$. The diagram reads as $\underline{\Sigma}_n(x)(i) = \underline{\Sigma}_0(X \bar{i} x) \quad (= \underline{\Sigma} \bar{i} x)$

(c) By (b)+(a), we have to get $\underline{\Sigma}_n(x)(i) = \underline{\Sigma}(X \bar{i} x)$

Easy: for $f: [m] \rightarrow [n]$ and $i \in m$; we get $\begin{cases} i \mapsto \underline{\Sigma}(X f \bar{i} x) \\ i \mapsto \underline{\Sigma}(X \bar{i} x) \end{cases} = (\text{since } f \circ \bar{i} = \bar{f})$

Boundaries $\partial\Delta^n$ and horns Λ_i^n

We give 3 definitions, from the more intuitive to the "less intuitive".

- As simplicial complexes:

$$\Delta^n = (\mathcal{V}, \mathcal{F}^*(\mathcal{V})) \quad (\mathcal{V} = \{v_0, \dots, v_n\})$$

$$\partial\Delta^n = (\mathcal{V}, \mathcal{F}^*(\mathcal{V}) / \{\mathcal{V}\})$$

$$\Lambda_k^n = \partial\Delta^n / \{\mathcal{V}_0, \dots, \mathcal{V}_k, \dots, \mathcal{V}_n\}$$

www.math.univ-paris13.fr/~valette/m1/cours/HomotopyTheory.pdf p.60

- Le bord $\partial\Delta^n$ du simplexe standard est égal au coégalisateur suivant:

$$\coprod_{0 \leq i < j \leq n} \Delta^{n-2} \xrightarrow{\quad} \coprod_{0 \leq l \leq n} \Delta^{n-1} \longrightarrow \partial\Delta^n,$$

où le morphisme en haut à gauche envoie la copie indiquée par $i < j$ sur celle indiquée par j via δ_i , le morphisme en bas à gauche envoie la copie indiquée par $i < j$ sur celle indiquée par i via δ_{j-1}

- Le k^{e} -cornet Λ_k^n de dimension n est égal au coégalisateur suivant:

$$\coprod_{0 \leq i < j \leq n} \Delta^{n-2} \xrightarrow{\quad} \coprod_{\substack{0 \leq l \leq n \\ l \neq k}} \Delta^{n-1} \longrightarrow \Lambda_k^n,$$

avec les mêmes morphismes.

See p. 27 for the real topological meaning!

www.math.uni-bonn.de/people/fhubeny/HigherSimpCats/Land-script.pdf p. 12

- (1) The boundary $\partial\Delta^n$ is the subsimplicial set of Δ^n whose k -simplices consist of the non-surjective maps $[k] \rightarrow [n]$.
- (2) For any subset $S \subseteq [n]$ the S -horn $\Lambda_S^n \subseteq \Delta^n$ consists of those k -simplices $f: [k] \rightarrow [n]$ where there exists a $i \in [n] \setminus S$ such that i is not in the image of f .

We then set $\Lambda_j^n = \bigcup_{\{j\}} \Lambda_j^n$

Exercise Show that these three definitions are equivalent.
Hint: for the first definition, use Sepr-to-Sect.

Spines

The third definition of horns / boundaries makes it clear that $\Lambda_i^n, \partial\Delta^n$ are subsimplicial sets (i.e. levelwise inclusions that are natural) of Δ^n .

$$\underline{\subset} \Delta^n$$

Another important family are the spines I^n . Again, the simplest definition is via simplicial complexes:

$$I^n = \left(\{v_0, \dots, v_n\}, \{\{v_0\}, \dots, \{v_n\}, \{v_0, v_1\}, \{v_1, v_2\}, \dots, \{v_{n-2}, v_n\}\} \right)$$

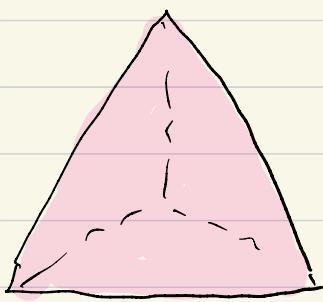
Exercise Show that spines are equivalently defined by

The spine $I^n \subseteq \Delta^n$ is given by those k -simplices $f: [k] \rightarrow [n]$ whose image is either of the form $\{j\}$ or $\{j, j+1\}$.

(from Land's notes p.1d)

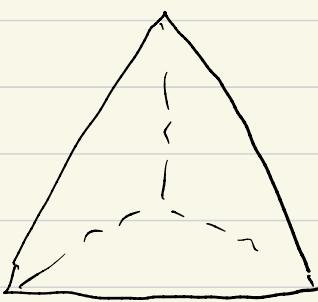
Give also a definition via a coequaliser.

In summary, and in pictures (dimension 3):



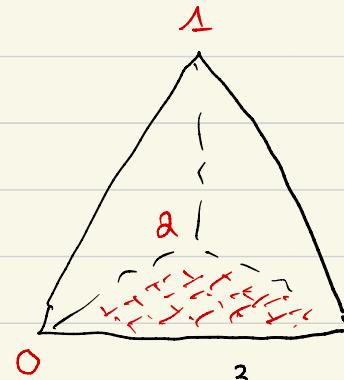
Δ^3
(everything)

of 3-ball



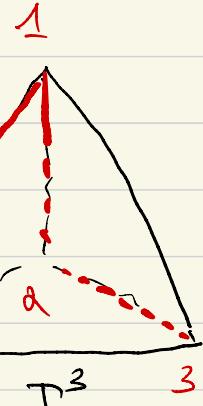
$\partial\Delta^3$
(fan Δ^2 , glued)

of 2-sphere



(three Δ^2 , glued)

all but the marked one, opposite to



I^3

Remark: We have

$$I^2 = \Lambda_1^2$$



Nerve of a category

There is a natural full and faithful functor $\underline{\subseteq} : \Delta \rightarrow \text{Cat}$

$$\text{Preorder } [0, \dots, n] \mapsto 0 \rightarrow 1 \rightarrow 2 \rightarrow \dots \rightarrow n$$

(pre)order as categories Recall the nerve functor from

Lecture 1: $N_{\subseteq} : \text{Cat} \rightarrow \widehat{\Delta}$

abbreviated as N^{\subseteq}

$$(NC)_n = \text{Cat}(\subseteq[n], \mathcal{C})$$

abbreviated as $\text{Cat}([n], \mathcal{C})$

$$\begin{array}{ccc} & f_0 & c_1 \\ c_0 & \nearrow & \searrow \\ & c_1 & c_2 \end{array} \dots$$

$f_{n-1} \downarrow c_n = \text{sequence of } n$
composable morphisms

Notation (f_0, \dots, f_{n-1})

$$\text{In particular } (NC)_0 = \text{Ob } \mathcal{C}$$

Face maps: $d_i(m \xrightarrow{f_i} \bigvee^{f_i} m) = m \xrightarrow{f_i \circ f_{i+1}} m$ (\circ is \sqcap = "inner")

$$d_0(m \xrightarrow{f_0} c_0) = m \quad d_n(m \xrightarrow{f_n} c_n) = m$$

Degeneracies $\delta_i(m \cdot m) = m \xrightarrow{id} m$

Exercise Prove that those are indeed the maps $\text{Cat}(\subseteq d^i, \mathcal{C}), \text{Cat}(\subseteq s^i, \mathcal{C})$
(Hint: spell out the action of \subseteq on d^i and s^i).

$$\text{(Visual hint: } \subseteq \left(\begin{matrix} 0 & 1 \\ 1 & 0 \\ 0 & 1 \\ 1 & 2 \end{matrix} \right) (0 \rightarrow 1) = 0 \rightarrow 1 \rightarrow 2 \text{)}$$

and therefore, $\underline{\subseteq} : \Delta \rightarrow \text{Cat}$ is dense
(cf. exercise p. 12 of lecture 1)

Exercise Show that $N : \text{Cat} \rightarrow \widehat{\Delta}$ is full and faithful.

(Hint: note that $NC_0 = \text{Ob } \mathcal{C}, NC_1 = \text{Mor } \mathcal{C}.$)

Since N is full and faithful, this raises the question of ...

Which simplicial sets are nerves of categories?

Nerve theorem

also called liftings

Land, p. 14-15

For a simplicial set X , the following three conditions are equivalent.

- (1) X has unique extensions for $\Lambda_j^n \rightarrow \Delta^n$ if $0 < j < n$.
- (2) X has unique extensions for $I^n \rightarrow \Delta^n$ for $n \geq 2$.
- (3) X is isomorphic to the nerve of a category.

$$\begin{array}{ccc} \Delta_j^n & \xrightarrow{\quad x \quad} & X \\ \downarrow & & \downarrow \\ \Delta^n & \xrightarrow{\exists! y} & \Delta^n \end{array} \quad \begin{array}{ccc} I^n & \rightarrow & X \\ \downarrow & & \downarrow \\ \Delta^n & \xrightarrow{\exists! z} & \Delta^n \end{array}$$

slightly

inner horns

(contrary = outer horns)

We shall make a \checkmark sharper statement and move it at the end of the lecture

wrong material introduced next:

- description of the left adjoint of N (from \mathbf{ASet} to \mathbf{Cat})
- skeleton and coskeleton

Notation $y \in \Delta_j^n$

Here is a preview of the sharper statement. We write

- $(1)^+$ for (1) where the inner restriction is removed for $n \geq 4$

Then we shall prove

$$(2) \Rightarrow (3) \Rightarrow (1)^+ \Rightarrow (1) \Rightarrow (2)$$

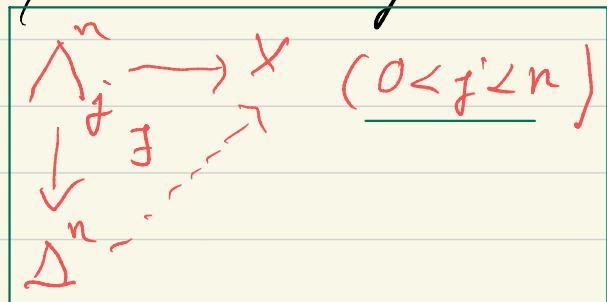
Exercise Show that if (1) holds for all horns ($0 \leq j \leq n$), then X is isomorphic to the nerve of a groupoid, and conversely.

A review of two approaches to higher categories

The two characterisations of the theorem on p. 74, when weakened, yield notions of weak category.

- A simplicial set satisfying (1) "minus uniqueness" is called a quasicategory. Explicitly: $(\Delta^{\text{op}}, 1)$

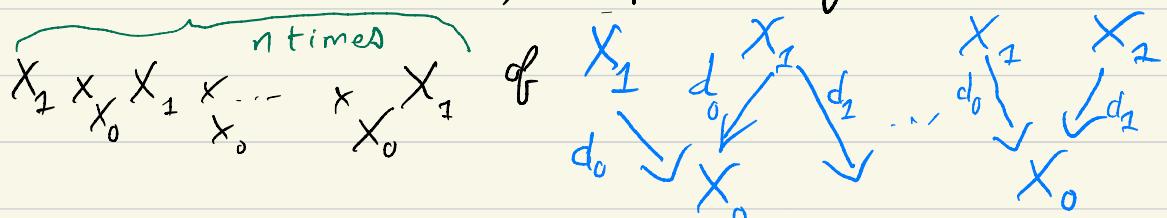
A quasicategory is a simplicial set X admitting liftings for all inner horns:



- Weakening (2) leads to another approach, due to Rozk. We first reformulate (2).

Exercise • Show that $\hat{\Delta}(I^n, X)$ is equivalently described as

the colimit



- and that (2) can be reformulated as (strict Segal condition)

The natural map $X_n \rightarrow X_1 \times_{X_0} \dots \times_{X_0} X_1$ is iso.

(in a suitable model category context)

By relaxing this iso to be only a $\xrightarrow{\sim}$ weak equivalence of simplicial sets (and hence taking X to be a bisimplicial set $\Delta^{\text{op}} \rightarrow \hat{\Delta}$) we get Rezk's notion of Segal space ...

$$(\Delta^{\text{op}} \times \Delta^{\text{op}} \xrightarrow{s} \Delta)$$

The other way around: from simplicial sets to categories

Consider

$$\Delta \xrightarrow{\subseteq} \text{Cat}$$

$$Y \downarrow \quad \text{Lan}_Y \subseteq$$

$$\Delta \leftarrow Y$$

Since Cat has all colimits, we can apply the formula.

$$\text{Lan}_Y \subseteq X = \underset{(n, x \in X_n)}{\text{colim}} \subseteq [n]$$

or "homotopy category"

also written $\text{h}X$

"realisation of X as a category"

It turns out that $\text{h}X$ has an alternative description, that we can synthesise: A cocone λ for the functor

$$(n, x) \mapsto [n] : \text{el}(X) \rightarrow \text{Cat}$$

with vertex C provides, in particular

$$\begin{array}{ccc} C & & \\ \nearrow & \searrow & \\ \lambda_{n,x} & & \text{h}X? \\ [n] & \nearrow x_{n,x} & \end{array}$$

- a map $\lambda_0 : X_0 \rightarrow \text{Ob } C$
 $\lambda_{0,-}$
- a map $\lambda_1 : X_1 \rightarrow \text{Mor } C$

- such that (naturality of λ) $\lambda_1(a \xrightarrow{a} a) = \text{id}_{\lambda_0(a)}(a)$ ($\forall a \in X_0$)

It is a good starting point for $\text{h}X$!

$$\text{Ob } (\text{h}X) = X_0$$

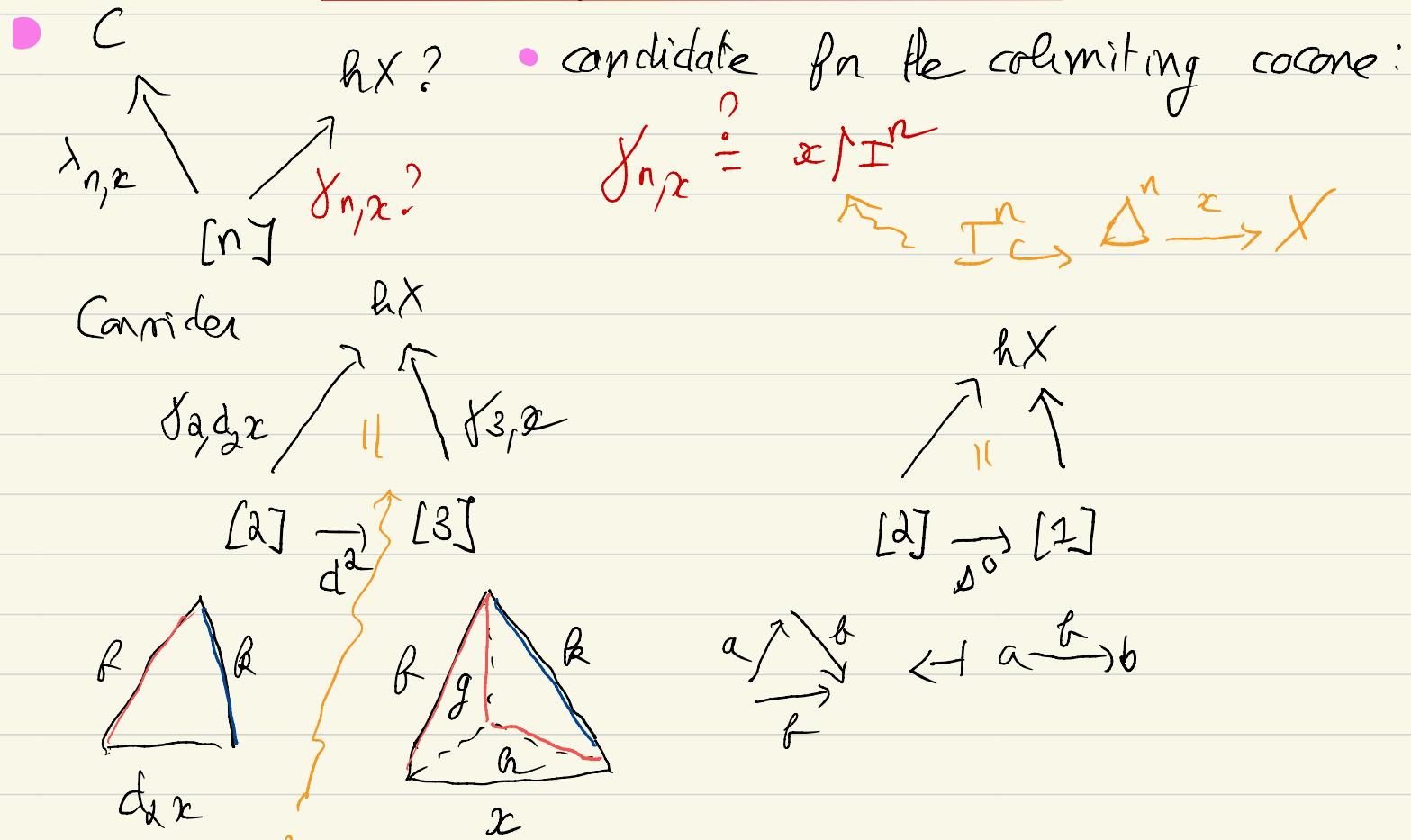
$$\text{Mor } (\text{h}X) = X_1$$

$$\forall a, \text{id}_a = a \xrightarrow{a} a$$

too naive!

- we need to define a colimiting cone γ
- we need to compose morphisms!

Synthesising hX , continued



Since $d^2(1 \rightarrow 2) = 1 + 2 + 3$
 $R = \text{hog}$ is forced in hX

Since $D^0(0 \rightarrow 1 \rightarrow 2) = 0 \rightarrow 1$
 $f \circ id_a = f$ is forced in hX

• Dealing with compositions: add them (as) greedily (as possible)

$$\text{Mor}(hX) \stackrel{?}{=} \bigsqcup_n \text{Spt}(I^n, X)$$

formal compositions of
composable 1-simplices

Combining • and • :



$$\text{Mor}(hX) = \bigsqcup_n \text{Spt}(I^n, X) / m \xrightarrow{b} \vee \text{or } m \xrightarrow{h} m,$$

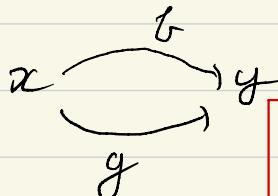
$$m \xrightarrow{b} \vee \xrightarrow{b} w = m \xrightarrow{b} w \quad m \xrightarrow{a} \vee \xrightarrow{a} w = m \xrightarrow{a} w$$

Exercise Check that $f_{n,x} = [x \sqcap I^n]$ is indeed a colimiting cocone.
 equivalence class \rightsquigarrow

The Boardman-Vogt construction

If X satisfies some conditions, then there is yet another description of πX . We define a relation on parallel 1-simplices, i.e.

We write $f \sim g$



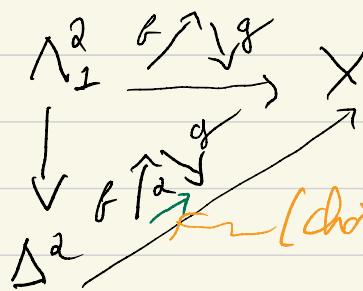
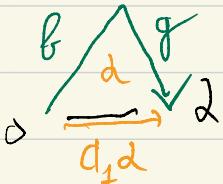
if $f_2 \sim g_2$ st. $x \xrightarrow{f} y$

Theorem If X has the (non necessarily unique) lifting property for Λ_2^2 and Δ_2^3 and Λ_2^3 (= all inner 2 and 3-horns), then \sim is an equivalence, and the following defines a category $\pi(X)$ which is isomorphic to πX :

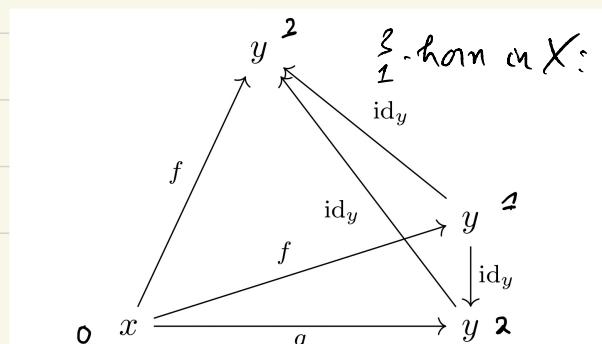
- $\text{Ob } \pi(X) = X_0$
- $\pi(X)(x, y) = \{f \in X_1 \mid d_1 f = x \text{ and } d_0 f = y\}/\sim$
- Moreover, we have $f \sim g$ iff $\begin{array}{c} f \\ \downarrow \\ x \end{array} \sim \begin{array}{c} g \\ \downarrow \\ y \end{array}$

Exercise Prove this theorem. Hints: • candidates for composition

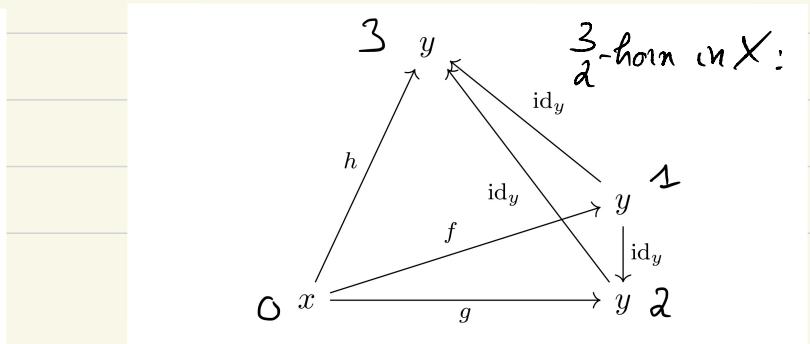
are given by Λ_1^2 liftings:



- For symmetry, build this

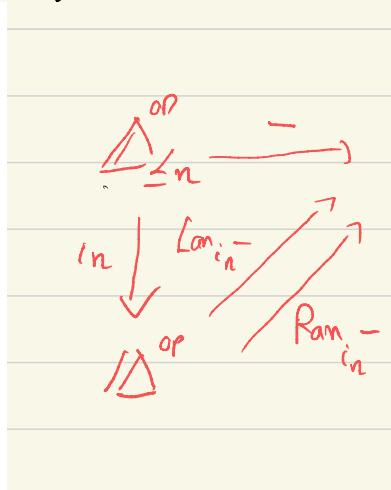
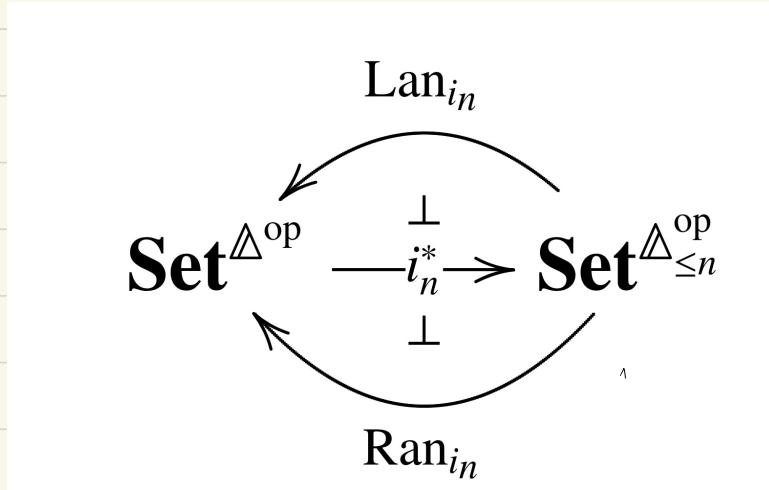


- For transitivity, build this



Skeleton and coskeleton

Let $\Delta_{\leq n}$ be the full subcategory of Δ on objects $\{i\}$ ($i \leq n$).



We write $\text{Ak}_n = \text{Lan}_{i_n} \circ i_n^*$

$\text{coker}_n = \text{Ran}_{i_n} \circ i_n^*$

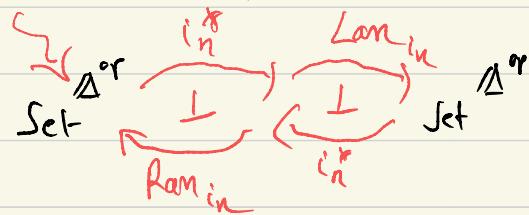
comonad

monad

on $\text{Set}^{\Delta^{op}}$

$$\begin{cases} & \\ & \end{cases} = \Delta$$

We have $\text{Ak}_n \dashv \text{coker}_n$ (as functors)



more precisely via
 ϵ , the counit of
the comonad

We say that X is n -skeletal if $X \cong \text{Ak}_n X$

X is n -coskeletal if $X \cong \text{coker}_n X$

η , unit iso

Skeleton in details

Explicitly, we have $\Delta^{\text{op}} \xrightarrow{i_n} \Delta^{\text{op}} \xrightarrow{\text{Set}}$
 in full and faithful \sim $i_n \downarrow$ " Δ^{op} $\xrightarrow{\text{Sk}_n X}$ left Kan ext.

$$(\text{Sk}_n X)_m = \bigcup_{i \leq n} \Delta^{\text{op}}([i], [m]) \cdot X_i = \{ (i, f \in \Delta([m], [i]), x \in X_i) \} / \text{gluing conditions}$$

We note that all m simplices in $\text{Sk}_n X$, for $m > n$, are degenerate (cf. first exercise on p. 47!)

Hence $\text{Sk}_n(X)$ has

- the same simplices as X in dimensions $\leq n$,
- and, above n , has only degeneracies of those.

Exercise Show that $\partial \Delta^n$ (as defined earlier) is $\text{Sk}_{n-1} \Delta^n$.

Exercise Show that for arbitrary X , $hX \simeq h(\text{Sk}_2 X)$.

Note that $i_d: X_{i_n} \rightarrow X_{i_n}$ induces a natural transformation $\text{Sk}_n X \rightarrow X$

Exercise Spell out this natural transformation, and that under the identifications above, it is the inclusion $\text{Sk}_n(X) \subseteq X$.

Cookelton in details

In contrast, $\text{cof}_n(X)$ has many m -simplices for $m > n$.

We have :

$$(\text{cof}_m X)_m \simeq \widehat{\Delta}(\Delta^m, \text{cof}_n X) \simeq \widehat{\Delta}(\partial \text{f}_n \Delta^m, X)$$

Yoneda rk-cof

In English, this gives the following description

called an n -shell of $(\text{cof}_n X)_m$ for $m > n$: each consistent decoration of the n -faces of Δ^m by n -simplices of X gives rise to a unique simplex in $(\text{cof}_n X)_m$ (the filler)

Exercise • Spell out the meaning of "consistent" here.

• Note that if all simplices in the shell are non-degenerate, then the unique filler is non-degenerate.

Exercise Synthesise the same description of cof_n from its definition by the end formula

$$(\text{cof}_n X)_m = \int_{\substack{* \\ i \leq n}} X_i \Delta^{([i], [m])}$$

(hint: in this formula, we can replace = restrict

$\Delta^{([i], [m])}$ by the injective morphisms $[i] \rightarrow [m]$).

Preparations for the nerve theorem

Exercise We have that $\rho k_2(\Lambda_j^n) \rightarrow \rho k_2(\Delta^n)$ is iso $\forall n \geq 4, 0 \leq j \leq n$
 (Hint all (iterated) faces of the missing face are (iterated) faces
 of another face in Δ^n .)

Lemma We have $h(\Lambda_1^3) = h(\partial\Delta^3) = h(\Lambda_2^3)$.

Proof $\partial\Delta^3 =$



as sequences of composable arrows, and they are equivalent thanks to
 the presence of the 0-face and of the 3-face in $\partial\Delta^3$

The same is true for Λ_1^3 and Λ_2^3 (but not for Λ_0^3, Λ_3^3 !)

Lemma Nerves of categories are 2-cokeletal.

Proof A 2^n -shell of NC contains in particular decorations

$$\underbrace{f_i}_{\stackrel{i \neq 2}{f_{i \circ f_i}}} \quad \underbrace{f'_i}_{\stackrel{i \neq 2}{f_i}} \quad \text{for all } i \quad \underbrace{f''_i}_{\stackrel{i \neq 2}{f_i}} \quad \text{Compatibility forces } f_i = f_{i \circ f_i}.$$

This already determines a candidate for the filling: $(f_0, \dots, f_{n-1}) \in NC_n$.
 This candidate is not only unique, but it also works, i.e. for every
 $\{i < j < k\} \subseteq [n]$ we have that $d_{\{i,j,k\}}(f_0, \dots, f_{n-1})$ is the face
 prescribed by the original 2-shell (shown by compatibility).

Exercise Show that the converse does not hold.

(Hint: several 2-simplices could have the same restriction to Δ^2 .)

Nerve theorem for simplicial sets (full statement)

Theorem TFAE for a simplicial set:

(1) X has unique liftings for all over n -horns ($n \geq 2$)

(1)⁺ X has unique liftings for all horns except for $\Delta_0^2, \Delta_2^2, \Delta_0^3, \Delta_2^3$

(2) X has unique liftings for all n -spines ($n \geq 2$)

(3) X is isomorphic to the nerve of a category

Proof We prove (2) \Rightarrow (3), (3) \Rightarrow (1)⁺, and (1) \Rightarrow (2)

- (2) \Rightarrow (3) We can form a category C by setting

$C = X_0$ and $C = X_1$, $\text{id}(c) = p_0 c$, $g \circ f$ given by the unique extension $b \nearrow g$ of $b \nearrow f: I^2 \rightarrow X$. (third side of the)

The monoid laws are ensured by uniqueness of the extensions of

$$\begin{array}{ccc} a & \xrightarrow{f} & b \\ a \downarrow & \nearrow g & \\ a & & \end{array} \quad \begin{array}{ccc} & b & \\ & \nearrow f & \searrow g \\ b & & b \end{array} \quad \begin{array}{ccc} f & g & h \\ \swarrow & \uparrow & \searrow \\ \text{both } f \text{ id and } f \\ \text{both } (hg)f \text{ and } h(gf) \end{array}$$

We then see that $X \simeq NC$ by (2):



$$X_n = \widehat{\Delta}(\Delta^n, X) = \widehat{\Delta}(I^n, X) = NC_n$$

Yoneda (2) by def. of C

• $(3) \Rightarrow (1)^+$

- For $n=2$, we have directly $\widehat{\Delta}(\Lambda_2^2, NC) \simeq NC_2 (\simeq \widehat{\Delta}(\Delta^2, NC))$ by definition of NC .

- For $n > 2$, we use some "categorical abstract nonsense".
We know that NC is 2 -cokeletal (second lemma p.23).
If Y is 2 -cokeletal and X is arbitrary, we have bijections

$$\begin{array}{c} \widehat{\Delta}(sk_2(X), Y) \\ \xrightarrow{\quad \text{adjunction} \quad} \\ \widehat{\Delta}(X, coh_2 Y) \\ \xrightarrow{\quad \text{cokeletal} \quad} \\ \widehat{\Delta}(X, Y) \end{array}$$

Our goal is to prove (for the specified forms) :

$$\frac{\widehat{\Delta}(\Lambda_k^n, NC)}{\widehat{\Delta}(\Delta^n, NC)} \stackrel{2\text{-cokeletality}}{\Leftrightarrow} \frac{\widehat{\Delta}(sk_2(\Lambda_k^n), NC)}{\widehat{\Delta}(sk_2(\Delta^n), NC)} \stackrel{n=1}{\Leftrightarrow} \frac{\text{Cat}(h(sk_2(\Lambda_k^n)), C)}{\text{Cat}(h(sk_2(\Delta^n)), C)}$$

and hence to prove $h(sk_2(\Lambda_k^n)) = h(sk_2(\Delta^n))$:

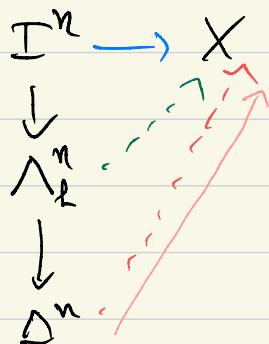
- for $n \geq 4$, we know (exercise p.23) $sk_2(\Lambda_k^n) \simeq sk_2(\Delta^n)$, for all forms, and we are a fortiori done.
- For $n=3$, we have $sk_2(\Lambda_k^3) = \Lambda_k^3$, hence our goal reformulates as $h(\Lambda_k^3) = h(sk_2(\Delta^3)) = h(\partial \Delta^3)$ for $k=1, 2$,

exercise p.21

- first lemma on p.23

• $(1) \Rightarrow (2)$

- We use induction on n . For $n=2$, this is tautological, as $\lambda^2 = \pm^2$
- Inductive case. It is enough to prove that every n -spine in X has a unique extension to some fixed chosen inner n -horn in X , because we can then use (1) for the further lifting to Δ^n .



Overall uniqueness of \rightarrow wrt to \rightarrow is deduced as follows: if \rightarrow is another candidate, then $\downarrow = \rightarrow$ by uniqueness, (claimed)
and $\rightarrow = \rightarrow$ by uniqueness (assumed)

We start with a decoration (f_0, \dots, f_{n-2}) of the spine of Δ^n . By assumption for $n=2$, we have unique compositions $f_i \circ f_j$ for all i , so that we can build $n+1$ ($n-1$)-spines in X :

$$(f_1 \cdots f_{n-2}), (f_1 \circ f_0, f_2, \dots, f_{n-2}), \dots, (f_0, \dots, f_{n-2} \circ f_{n-2}), (f_0, \dots, f_{n-2})$$

$x_0 \quad x_1 \quad \quad \quad x_{n-1} \quad x_n$

By induction for $n-1$, we get $n+1$ simplices $x_i \in X_{n-1}$ having those spines

We are left to show to show:

If $\{i < j\} \subseteq [n] \setminus \{i_0\}$ for some chosen $0 < i_0 < n$

$$(x_i)_{j-1} = (x_j)_i$$

By assumption for $n=2$, it is enough to look at the spines of (x_{ij}) and $(x_j)_i$, which are again obtained from $\text{spine}(x_i)$ and $\text{spine}(x_j)$ by the same process as above

Exercise Before examining general n , work out $n=3$.
(next page)

(1) \Rightarrow (2) continued

$$(f_1 \cdots f_{n-1}), (f_1 \circ f_0, f_2, \dots, f_{n-2}), \dots, (f_0, \dots, f_{n-2} \circ f_{n-1}), (f_0, \dots, f_{n-2})$$

$x_0 \quad x_1 \quad \quad \quad x_{n-1} \quad \quad \quad x_n$

- $i=0, j=n$ We have $(x_0)_{n-1} = (f_1, \dots, f_{n-2}) = (x_n)_0$
- $i=0, j=1$ $(x_0)_0 = (f_2, \dots, f_{n-2}) = (x_1)_0$
- $i=0, 1 < j < n$ We have $(x_0)_{j-1} = (f_1, \dots, f_j \circ f_{j+1}, \dots, f_{n-2}) = (x_j)_0$
- $0 < i < j < n, j-i > 1$ We have

$$(x_i)_{j-1} = (f_0, \dots, (f_i \circ f_{i-1}), \dots, (f_j \circ f_{j-1}), \dots, f_{n-2}) = (x_j)_i$$
- $0 < i < n-1, j=i+1$ $(x_i)_i = (f_1, \dots, f_{i+2} \circ (f_i \circ f_{i-1}), \dots, f_{n-2})$

$$(x_{i+2})_i = (f_1, \dots, (f_{i+2} \circ f_i) \circ f_{i-1}, \dots, f_{n-2})$$

Here we distinguish two cases :

- $n > 3$. Then we can apply induction for $n=3$ to the 3-apme (f_{i-2}, f_i, f_{i+2}) and get a 3-apmper in X witnessing associativity, and then $(x_i)_{i+2} = (x_{i+2})_i$.
- $n=3$. Then we simply **avoid** this case by requiring our chosen horn to be an inner horn : indeed, in an inner n -horn, x_0 and x_n have to be there. If $n=3$, there is only one left face to be chosen to form an inner horn : x_1 and x_2 cannot be both in, i.e. $0 < i < j < n$ impossible.
- All cases with $j=n$ are symmetric to the cases with $i=0$.

Remark : For $n \geq 4$, it does not matter which faced horn we choose (we could as well have chosen an outer horn, except that we need the horn to be inner for applying our assumption (1)).

And topology in all this?

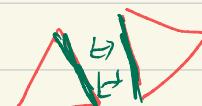
Replace Cat by Top, replace \subseteq by the functor

$\bar{\Delta}_{\text{top}} : \Delta \rightarrow \text{Top}$ defined by $\Delta_{\text{top}}^n = \text{convex hull of } \{(1, 0, \dots, 0), (0, 1, 0, \dots, 0), \dots, (0, \dots, 0, 1)\} \text{ in } \mathbb{R}^{n+1}$.

Exercise Describe the action of $\bar{\Delta}_{\text{top}}$ on morphisms

Then we have

$$\begin{array}{ccc} \Delta & \xrightarrow{\quad \text{Top} \quad} & \Delta \\ \downarrow \text{written } I-1 & \nearrow \text{long } \bar{\Delta}_{\text{top}} & \downarrow N \\ \widehat{\Delta} & \xrightarrow{\quad \text{S} \quad} & \end{array}$$



Instantiating the definitions, we have

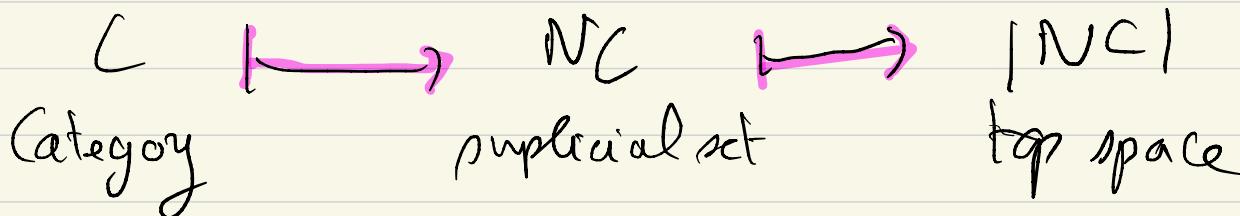
$(SA)_n = \text{Top}(\Delta_{\text{top}}^n, A)$ (the singular complex of A)

$|X| = \int^n X_n \times \Delta_{\text{top}}^n$ (the geometric realization of X)

Exercise Make sense of f^* and f_* in the following formula, and

prove it:

$$|X| = \left(\coprod_{n \in \Delta} X_n \times \Delta_{\text{top}}^n \right) / ((f^*(x), t) \sim (x, f_*(t)))$$



$$\begin{array}{ccc} \Delta & \xrightarrow{\quad \text{Cat} \quad} & \\ \downarrow & \nearrow \text{I-1} & \\ \widehat{\Delta} & \xrightarrow{\quad \text{S} \quad} & \end{array}$$

$$\begin{array}{ccc} \Delta & \xrightarrow{\quad \text{Top} \quad} & \\ \downarrow & \nearrow \text{I-1} & \\ \widehat{\Delta} & \xrightarrow{\quad \text{S}(\text{sing}) \quad} & \end{array}$$